

We did a var ~~dist~~ ~~approx~~ approx with a factored dist — BUT we could do others: ①

eg. Q: a free structured dist?

$$E_Q = -E_Q \log P \rightarrow H_Q$$

1) We can compute  $H_Q$

recall tree:

$$q(x_1 \dots x_N) = \left[ \prod_{i \in V} q_i \right] \left[ \prod_{i \in E} \frac{q_{ij}}{q_i q_j} \right]$$

$$= \frac{\prod_{i \in E} q_{ij}}{\prod_{i \in V} q_i^{(d_i-1)}}$$

degree of edge

$$\text{So } H_Q = - \sum_{\text{values of pairs}} \left[ \sum_{i \in E} q_{ij} \log q_{ij} \right] + \sum_{\text{values}} \left[ (d_i-1) \sum_{i \in V} q_i \log q_i \right]$$

tractable:

We can also compute

$$-E_Q \log P$$

recall:

$$P(H|x) = \frac{1}{Z} \exp \left[ - \sum_{ij} \theta_{ij}(H_i, H_j) - \sum_i \theta_i(H_i) \right]$$

$$- \log P = \log Z + \sum_{ij} \theta_{ij}(H_i, H_j) + \sum_i \theta_i(H_i)$$

$\uparrow$                        $\uparrow$                        $\uparrow$   
*x<sub>i</sub>*                      in other hand w  
 notes ; sorry !                      vars ;

constant - not a prob, cause we're minimizing!

So

$$-E_Q \log P = \log Z + \sum_{\text{values of pairs}} \left[ \sum_{ij \in \mathcal{Z}} q_{ij} t_{ij} \right] + \sum_{\text{values}} \left[ (d_i - 1) \sum_{i \in \mathcal{V}} q_i t_i \right]$$

Here I used the change of var

$$t_{ij} = \theta_{ij} + \theta_i + \theta_j$$

$$t_i = \theta_i$$

to get an expression that looks like entropy.

Now: we want to min

$$E_Q \log Q - E_Q \log P$$

for  $Q$  some fixed tree. (which we chose)

Recall the  $q_{ij}$ ,  $q_i$ ,  $q_j$  are marginals

so we must

$$\min \sum_{\text{values}} \left[ \sum_{ij \in E} q_{ij} [\log q_{ij} + t_{ij}] \right] + \sum_{\text{values}} \left[ \sum_i (d_i - 1) \cdot q_i [\log q_i + t_i] \right]$$

$$\text{s.t.} \quad \sum_i q_{ij} = q_j \quad ; \quad \sum_j q_{ij} = q_i \quad ; \quad \sum q_i = 1$$

(cause these are marginals)

write

$\lambda_{\epsilon_i}$  for LM's asso. with  
 $\lambda_{\epsilon_j}$  . . . . .  
 $\lambda_{v_i}$  LM . . . . .

$$\sum_j q_{ij} = q_i$$

$$\sum_i q_{ij} = q_j$$

$$\sum_i q_i = 1$$

(Notice  $\lambda_{\epsilon_i}$  is a vector  
 $\lambda_{v_i}$  scalar)

write lagrangian  $L$ .

at stationary point

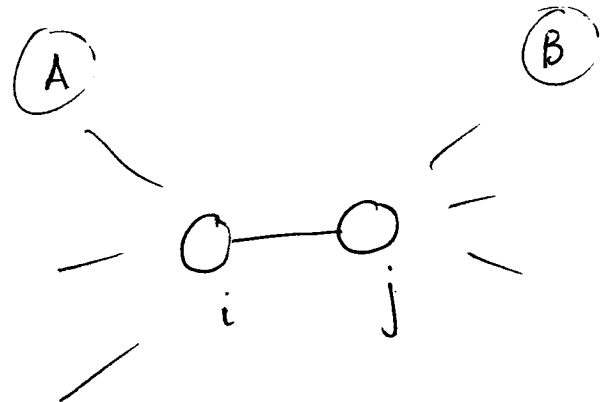
$$\left[ \frac{\partial L}{\partial q_{ij}} \right]_{uv} = 0 = \left[ \log q_{ij} \right]_{uv} + 1 + \left[ \tilde{z}_{ij} \right]_{uv} + \sum_{\text{all incoming edges to } i} \left[ \lambda_{\epsilon_i} \right]_u + \sum_{\text{all incoming edges to } j} \left[ \lambda_{\epsilon_j} \right]_v$$

$\left[ \frac{\partial L}{\partial q_{ij}} \right]_{uv}$   $\rightarrow$   $u, v$ 'th entry in table  
 $\left[ \log q_{ij} \right]_{uv}$   $\rightarrow$  ~~table~~  
 $\left[ \tilde{z}_{ij} \right]_{uv}$   $\rightarrow$   $\tilde{z}_{ij}$   
 $\sum_{\text{all incoming edges to } i} \left[ \lambda_{\epsilon_i} \right]_u$   $\rightarrow$  all incoming edges to  $i$   
 $\sum_{\text{all incoming edges to } j} \left[ \lambda_{\epsilon_j} \right]_v$   $\rightarrow$  all incoming edges to  $j$

so

$$\left[ q_{ij} \right]_{uv} \propto \left[ e^{-\tilde{z}_{ij}} \right]_{uv} \cdot \left[ e^{\sum \lambda_{\epsilon_i}} \right]_u \cdot \left[ e^{\sum \lambda_{\epsilon_j}} \right]_v$$

compare with B.P. eqns



$$[q_{ij}]_{uv} \propto [\psi_{ij} \phi_i \phi_j]_{uv} \cdot \left[ \prod_{\text{all inc to } i} M_{ia} \right]_u \cdot \left[ \prod_{\text{all inc to } j} M_{jb} \right]_v$$

Conclusion

- Messages = log  $\sum M$ 's

Two outcomes :

- 1) We can fit a variational model of a single tree (MP as above)
- 2) Hoopy BP "like" fitting var m of tree without worrying about tree

Now what is happening in terms of M.P. ? ⑥

→ fix a tree.

→ Interpret MP  $\equiv$  convex hull of all states that can arise in this repn of G.M.

→ ~~we must have that~~

Call the polytope that satisfies

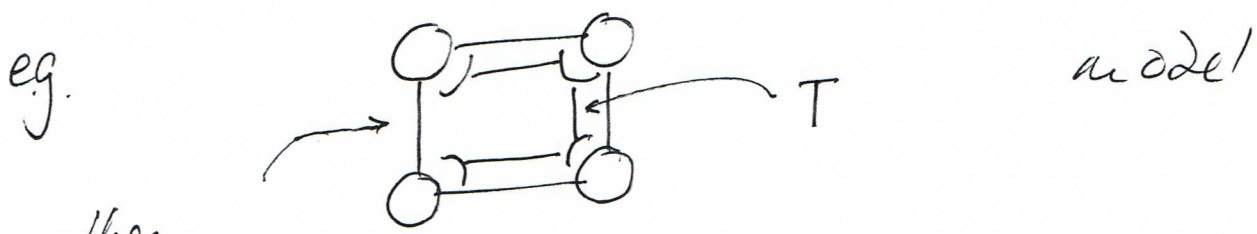
$$\sum_j q_{ij} = q_i \quad ; \quad \sum_i q_{ij} = q_j \quad ; \quad \sum_i q_i = 1$$

the Local Polytope =  $L_p$ .

→ Notice

$$MP \subset L_p.$$

Notice that for choice of tree  $T$ ,  
a lot of  $q_{ij} = q_i \cdot q_j$  (cause the  
vars are ~~indep~~ cond indep given parents



then  
in  $T$ , these  
two have  $q_{ij} = q_i \cdot q_j$

so we are finding the  $q_{ij}, q_i, q_j$   
that

- are in LP.
- meet indep constraints implied by  $T$
- minimize  $-E_Q \log P + E_Q \log Q$

(then ~~st~~ extract info from  $q$ ).

# Extracting info from Q.

- if we're lucky,  $q_{ij}$  are integer.  
(might be a vert ~~of~~  $M_p!$ ).
- nothing to do
- else, it's a tree; → max product

## Idea:

rather than

$$\min -E_Q \log P + E_Q \log Q$$

for  $Q$  a tree,

do it for  $Q \in LP$ .

→ How do we get  $E_Q \log Q$ ?



Here is one strategy

- drop the tree
- ~~compare~~ fit  $q$  by using the expression for  $E_Q \log Q$  that came from tree

$$E_Q \log Q \approx - \sum_{\text{edges}} \sum_{\substack{\text{values} \\ \text{of pairs}}} [q_{ij} \log q_{ij}] + \sum_{\text{verts}} \sum_{\text{values}} [(d_i - 1) \frac{q_i}{d_i} \log \frac{q_i}{d_i}]$$



⇒ notice I flipped order  
 =  $H_Q^?$

This isn't the true exp. for  $E_Q \log Q$ , but  
 its easy to eval.

↳ loopy b.p. =  $\min E_Q \log P - H_Q^?$   
 st.  $Q \in LP$ .

Notice the form of the costfunction

$$E_Q \log P \sim H^?$$

↑  
linear in Q.

↑  
Some property of Q that approx  
entropy.

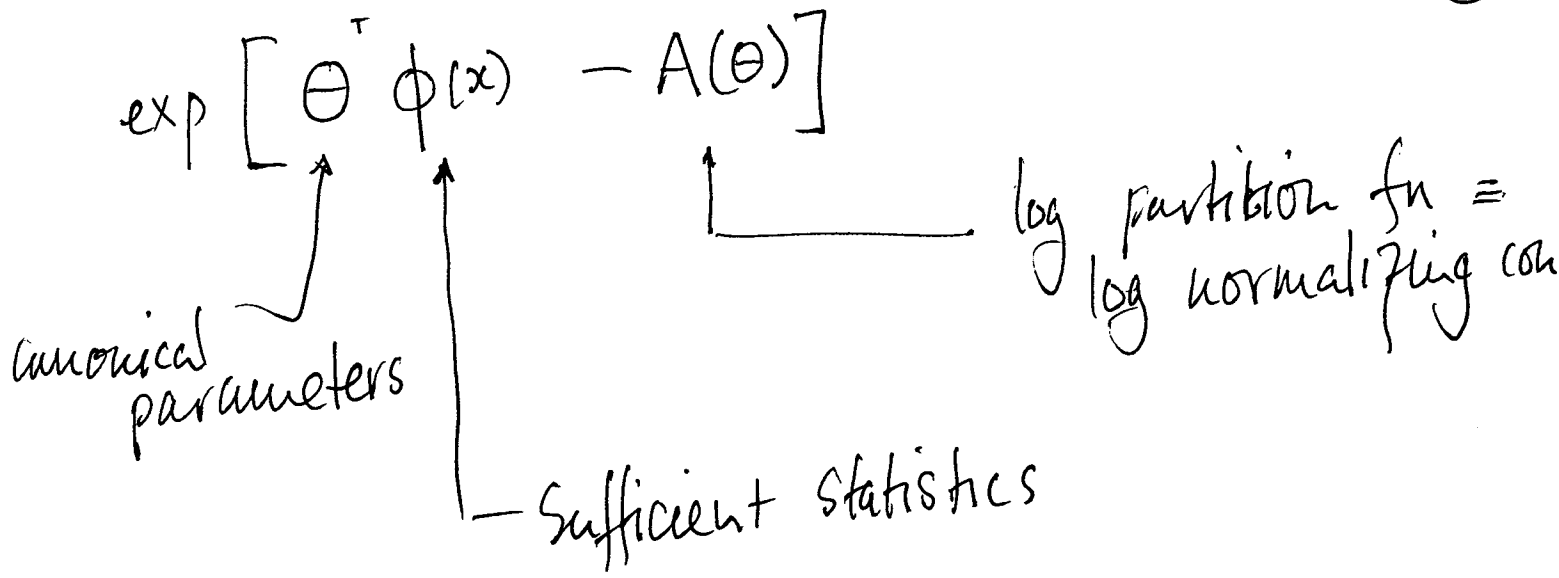
Notice also that we're identifying points in  $\mathcal{P}$  (or  $\mathcal{MP}$ ) with probability distributions. It turns out that we can formalize this

The exponential family:

any p.d. that is written as

$$P(x) = \exp \left[ \theta^T \phi(x) - A(\theta) \right]$$

(for our purposes - other possible)



We will confine attention to case where:

- $\phi(x)$  are linearly independent (no real issue here, just creates a lot of if's, and's, but)
- $\theta$  is such that

$$A(\theta) = \log \int \exp[\theta^\top \phi(x)] < \infty$$

& ~~that~~

Examples:

1D Normal Dist:

$$\exp \left[ \cancel{(\alpha, \beta)} \begin{pmatrix} \alpha \\ \beta \end{pmatrix}^T (x^2, x) - A(\Theta) \right]$$

here  $\alpha < 0$  ;  $\text{std} = \frac{-1}{2\alpha}$ .

$$\text{mean} = \left(\frac{1}{2\alpha}\right)^2 \cdot \beta$$

Multi D ND:

follows easily.

Poisson dist:

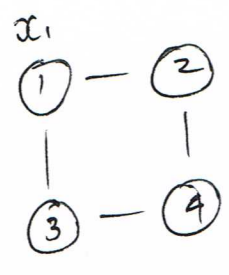
recall this is dist on ~~the~~ non-neg integers  
 ← rate/intensity.

$$P(k) = \lambda^k \frac{e^{-\lambda}}{k!}$$

$$p(k) = \exp [\alpha \cdot k - A(\alpha)]$$

$$\alpha = \log \lambda, \text{ etc.}$$

Discrete MRF :



etc.  $\mathbb{1}_{x_i}$  is 1-hot vector  
 $\mathbb{1}_{x_i, x_j}$  is ~~the~~ 1-hot table, straightened into vector

$$p(x) = \exp \left[ \Theta^T \begin{bmatrix} \mathbb{1}(x_i) \\ \vdots \\ \mathbb{1}(x_i, x_j) \\ \vdots \end{bmatrix} - A(\Theta) \right]$$

$\Lambda(\theta)$  is extremely interesting

$$\begin{aligned} \nabla_{\theta} A &= \nabla_{\theta} \left[ \log \int e^{\theta^T \phi} dx \right] \\ &= \frac{1}{\int e^{\theta^T \phi} dx} \cdot \int \phi e^{\theta^T \phi} dx \\ &= E_{\rho}[\phi] \end{aligned}$$

recall - we've seen something like this before when talking about max-likelihood = max entropy.

Now assume we have some  $\phi$  (likely indicator fns in our case).

We can define  $\Lambda : \theta \rightarrow M$  ← marginal polytope

$$\Lambda(\theta) = E_{\theta}[\phi]$$

↑ this is in  $M$ , cause  $M$  is all possible expectations of

Then:

$\Delta$  is 1-1 (assuming  $\phi$  are linearly indep).

(proof in Wainwright - mildly technical).

Now we want to consider dual of  $A(\theta)$ .

$$A^*(\mu) = \sup_{\theta \in \Theta} [\langle \mu, \theta \rangle - A(\theta)]$$

Note this is a function of  $\mu$ .  
Known as a conjugate dual

Why is  $\Lambda$  1-1?

15a

(Sketch of proof - details in mainwright)

$A(\theta)$  is convex  
we must show for any  $\mu$ , there  
is some  $\theta$  st  $E_{p(x;\theta)}[\phi] = \mu$ .

BUT  $E_{p(x;\theta)}[\phi] = \nabla_{\theta} A(\theta)$ .

under very mild conditions, map  
 $x \rightarrow \frac{df}{dx}$  is 1-1 for  $x$  convex

- proof by drawing!



$A(\theta)$  is a convex function (16)  
of  $\theta$ .

recall  $\frac{\partial A}{\partial \theta_i} = e^{-A} \cdot \int e^{\theta^T \varphi} \cdot \varphi_i dx$

so  $\frac{\partial^2 A}{\partial \theta_i \partial \theta_j} = e^{-A} \cdot \int e^{\theta^T \varphi} \varphi_i \varphi_j dx$   
 $- \left[ e^{-A} \cdot \int e^{\theta^T \varphi} \varphi_j dx \right] \left[ \int e^{\theta^T \varphi} \varphi_i dx \right]$

$$= E_P[\varphi_i \varphi_j] - E_P[\varphi_i] E_P[\varphi_j]$$

$$= \text{cov}(\varphi_i, \varphi_j)$$

so  $H_A = \text{covmat}[\varphi]$

this is our positive definite conditions (under linearly indep.)

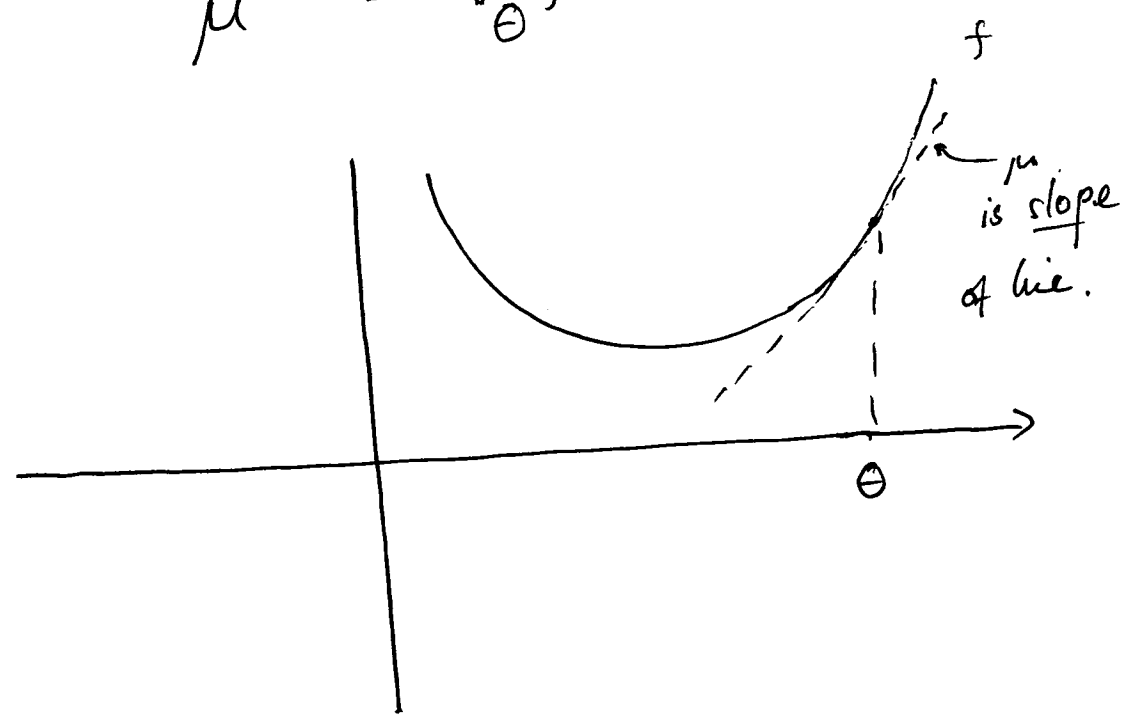
Now consider conjugate dual.

$$f^*(\mu) = \sup_{\theta} [\langle \mu, \theta \rangle - f(\theta)]$$

for convex  $f$ .

- assume  $f$  differentiable

then  $\mu = \nabla_{\theta} f$

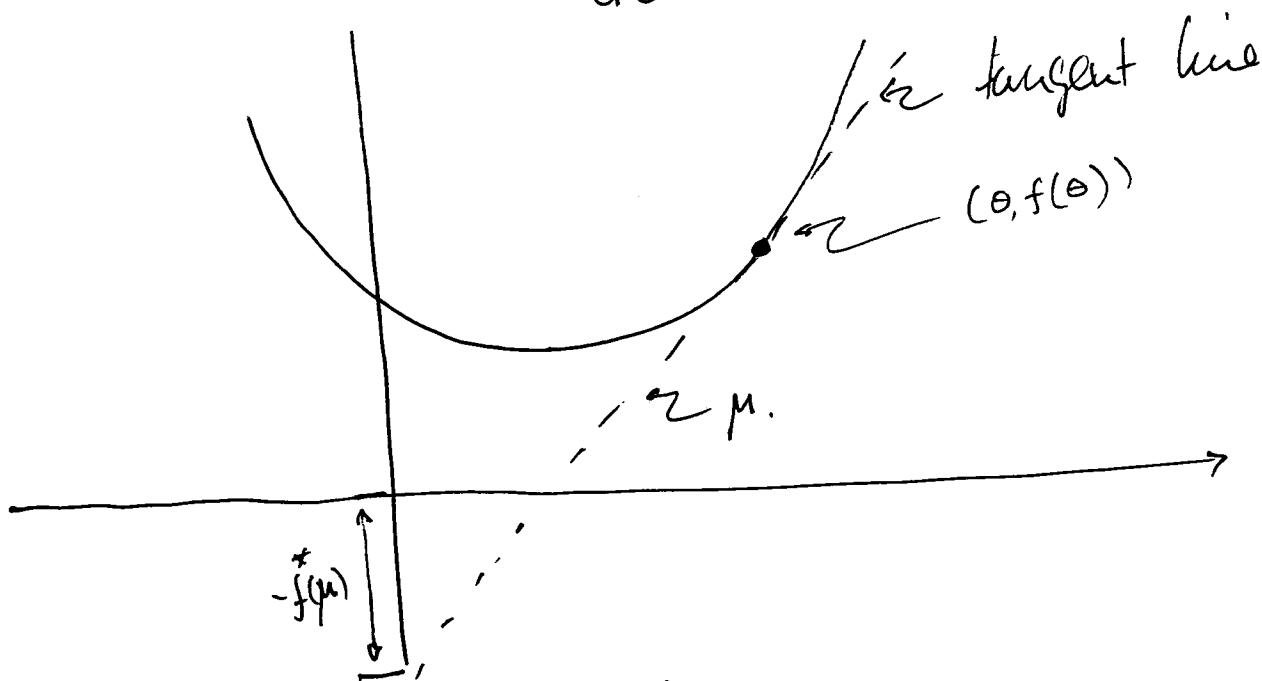


# Common visualization

(18)

- choose  $\mu \rightarrow$  what is  $f^*(\mu)$ ?
- consider  $\theta$  such that

$$\mu = \frac{df}{d\theta}$$



tangent line has slope  $\mu$ .  
passes through  $(\theta, f(\theta))$

$$\therefore y = \mu x + (f(\theta) - \mu \theta)$$

$\therefore$  at  $x = 0$ ,  $y = -f^*(\mu)$

for  $f$  convex,  $f^*$  is convex.

(19)

Show for  $f \in C^2$ , but generally true.

Proof:  $f^*(p) = \sup_x [px - f(x)]$

$f$  diff, convex so

$$p = \frac{df}{dx} \quad \text{at sup.}$$

$\frac{df}{dx}$  is a function, and ~~each~~ is 1-1 exists

so  $g$  s.t.  $g \circ \frac{df}{dx} = \text{Id}$

$$g(p) = x \quad \text{at sup}$$

so  $f^*(p) = p \cdot g(p) - f(g(p))$

$$\frac{df^*}{dp} = p \frac{dg}{dp} + g - f' \cdot \frac{dg}{dp} = g(p)$$

so  $\frac{d^2 f^*}{dp^2} = \frac{dg}{dp} = \frac{dx}{dp}$ ; but  $\frac{dp}{dx} = \frac{df}{dx^2}$

$$d^2 f^* = \frac{1}{\|f''\|} > 0$$

all this works in  $\mathbb{N}^D$  as well  
(ex: prove it!)

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Thm (Fenchel - Moreau).

$$f = (f^*)^*$$

iff

$f$  is proper, lower semi-continuous,  
and convex

OR

$$f \equiv \infty$$

OR

$$f \equiv -\infty$$

Now consider  $A(\theta) = \log Z(\theta)$ .  
for an exp. dist.

1)  $A^*(\mu) = \sup_{\theta} (\langle \theta, \mu \rangle - A(\theta))$   
is defined, convex.

2) for  $\mu \in \mathcal{M}$   
marginal polytope.

write  $\theta(\mu) = \Lambda^{-1}(\mu)$

then  $A^*(\mu) = -H(p(x; \theta(\mu)))$ .

Proof of 2 (sketch):

(22)

$$\Lambda^{-1}(\mu) = \theta \quad \text{such that}$$

$$E_{p(x; \theta)} [\phi(x)] = \mu$$

(by defn of  $\Lambda$ )

$$\text{But if } \mu = E_{p(x; \theta)} [\phi] = \nabla_{\theta} A(\theta)$$

then  $\theta$  is sup

$$\begin{aligned} \text{so } -H(p(x; \theta(\mu))) &= E_{p(x; \theta(\mu))} [\langle \theta, \phi(x) \rangle - A(\theta)] \\ &= \langle \theta, \mu \rangle - A(\theta) \\ &= A^*(\mu) \quad (\text{cause } \theta \text{ is sup}) \end{aligned}$$

③

$$A(\theta) = \sup_{\mu \in M} \{ \langle \theta, \mu \rangle - A^*(\mu) \}$$

(A is lower semicontinuous - see notes;  
then Fenchel-Moreau means

$$(A^*)^* = A$$

and  $(A^*)^*(\theta) = \sup_{\mu \in M} \{ \langle \theta, \mu \rangle - A^*(\mu) \}$

compare : with

$$E_Q = -E_Q \log p + E_Q \log q$$

which we minimize to build var model  
marginals of Q

$$E_Q \log p \rightarrow \langle \theta, \mu \rangle$$

↑  
params of



Now we have .

$$\sup_{\mu \in \mathcal{M}} \{ \langle \theta, \mu \rangle - A^*(\mu) \} \quad \leftarrow \textcircled{C}$$

is attained at

$$\mu = E_{p(x; \theta)} [\phi]$$

So solving  $\textcircled{C}$  gives

- log partition function
- Set of mean pars  
(for our purposes,)  
arg Max .

BUT

$\mathcal{M}$  is hard.

$A^*(\mu)$  is hard.

Now we can unify algs.

Mean field, single tree, etc

for any  $\mu \in M$ ,

$$A(\theta) \geq \langle \mu, \theta \rangle - A^*(\mu).$$

Now consider  $T \subset M$

↳ corresponding to models that are tractable  $\equiv$  can compute  $A^*(\mu)$

then solve

$$\sup_{\mu \in T} \{ \langle \theta, \mu \rangle - A^*(\mu) \} = A_{MF}(\theta)$$

↳ same as our exp but w - sign of max

must have  $A_{MF}(\theta) \leq A(\theta)$

loopy BP

1) recall for a tree structured model

$$H(q) = - \sum_{\substack{\text{# values} \\ \forall i \in \text{verts}}} \left[ \sum_{x \text{ values}} q_i(x_i) \log q_i(x_i) \right] \\ - \sum_{i,j \in \text{edges}} \left[ \sum_{x_i, x_j \text{ values}} q_{ij}(x_i, x_j) \cdot \log \left[ \frac{q_{ij}(x_i, x_j)}{q_i(x_i) q_j(x_j)} \right] \right]$$

2) approximate

$$A^*(\mu) \approx -H(\mu) \leftarrow \begin{array}{l} \text{computed using} \\ \text{tree expression} \\ \text{Bethe approx} \end{array}$$

3)  $\mathcal{L} \supset \mathcal{M}$

↑  
 local polytope,  
 consistency constraints  
 for pairwise marginals

4) Solve

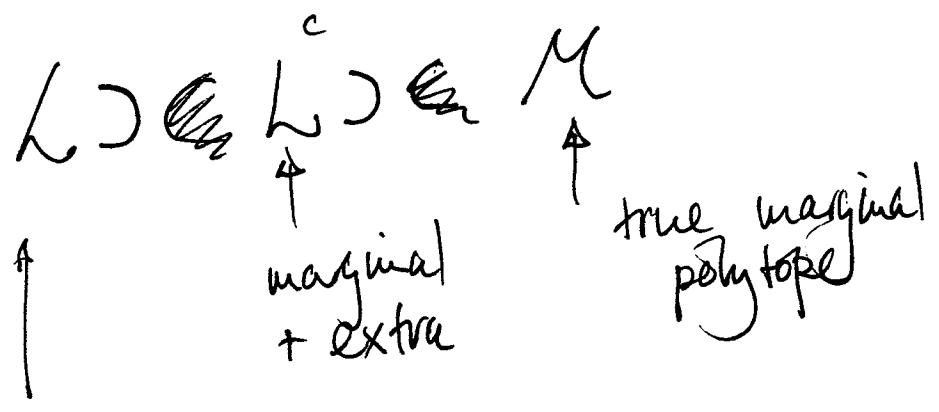
$$\sup_{\mu \in \mathcal{L}} \{ \langle \theta, \mu \rangle + H_B(\mu) \} = A_{LBP}(\theta)$$

we must have

$$A_{LBP}(\theta) \geq A(\theta)$$

But we can now explore other approximations:

- eg. insert constraints so that



defined by marginal constraints

here's one construction.

Assume some vector  $\mu$ , which might be in  $M$

Construct the matrix

$$M = \begin{bmatrix} 1 & \mu_1 & \mu_2 & \dots \\ \mu_1 & \mu_{11} & \mu_{12} & \dots \\ \mu_2 & \mu_{21} & \mu_{22} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

$\mu_1, \mu_2$   
means  
means for first var

$M$  is a covariance matrix  
so  $M \succeq 0$