

Interior point methods

①

problem

$$\min f(x)$$

$$\text{s.t. } g_i(x) \leq 0$$

assume that feasible set is not empty

Strategy: avoid the boundary, cause progress on the boundary may be slow

- hard to take large steps, cause active constraints might change
- overstep: - recovering might be hard

Instead, introduce a barrier function,

- ∞ at boundary
- large close to boundary

then solve; make barrier closer to boundary
repeat.

Generally, we engage only w/convex \mathbb{R}^n feasible sets

- More interesting than you think.
- CURRENTLY, only linear constraints (but I'll show others soon!)

Barrier

$$\phi(x) = -\sum_{i=1}^M \log(-g_i(x))$$

(signs depend on direction of ineq.!) ϕ is finite

Notice:

$$g_i(x) < 0$$

~~and~~

$$g_i(x) = 0$$

ϕ is $+\infty$

$$g_i(x) \text{ very neg}$$

~~and~~ ϕ_i is small

Now consider

$$\min_x f(x) + \frac{1}{t} \varphi(x)$$

at small t , solu is far in interior
big t , solu can approach boundary

Usually : phrase as

$$\min_x t f(x) + \varphi(x)$$

It is usual to have linear constraints
as well, so

$$\begin{aligned} \min_x & t f(x) + \varphi(x) \\ \text{st} & Ax = b \end{aligned}$$

Notice:

f convex, g_i convex means
 $t f + \varphi$ is convex.

Now consider $x^*(t)$ which is soln

$$A x^* = b$$

$$g_i(x^*) < 0$$

(cause it is soln!)

$$\nabla_x \mathcal{L} = t \nabla_x f(x^*) + \nabla_x \varphi(x^*) + \gamma^T A = 0$$

grad of lagrangian is zero at soln

$$\nabla_x \mathcal{L} = 0 = t \nabla_x f(x^*) + \sum_i \left[\frac{-1}{g_i(x^*)} \nabla_x g_i(x^*) \right] + \gamma^T A = 0$$

So:

$$\nabla_x f(x^*) + \sum_i \left[\frac{-1}{t g_i(x^*)} \right] \nabla_x g_i(x^*) + \frac{\nu^T}{t} A = 0 \quad (5)$$

write:

$$\lambda_i^* = \frac{-1}{t g_i(x^*)}, \quad \nu^* = \frac{\nu}{t}$$

then

$$\nabla_x f(x^*) + \sum_i \lambda_i^* \nabla_x g_i(x^*) + \nu^{*T} A = 0$$

$$\lambda_i^* \geq 0$$

but this is dual feasibility for the original problem, cause x^* is feasible

Furthermore, we have that

⑥

x^* minimizes

$$L(x, \lambda^*, \nu^*) = f + \sum_i \lambda_i^* g_i + \nu^{*T} (Ax - b)$$

So: the dual is ~~not~~ not infinite at x^* [(λ^*, ν^*) are dual feasible]

Further, we can evaluate the dual at (λ^*, ν^*)

$$g(\lambda^*, \nu^*) = \text{dual} \\ = \inf_x L(x, \lambda^*, \nu^*)$$

$$= L(x^*, \lambda^*, \nu^*) \quad [\text{cause } x^* \text{ minimizes}]$$

$$= f(x^*) + \sum_{i=1}^m \left[\frac{-1}{t} g_i(x^*) \right] + \nu^{*T} [Ax - b]$$

$$\uparrow \\ -\frac{m}{t}$$

$$\uparrow \\ 0$$

So: ⑦

- primal has value $f(x^*)$
- dual has value $f(x^*) - \frac{m}{t}$

• BUT dual is lower bound on value of the problem ($q(\lambda, \nu) \leq f(x)$).

• so soln is trapped in range $[f(x^*) - \frac{m}{t}, f(x^*)]$

Algorithm:

- Start with x a feasible point,
 t small ($t=1$, say).
 $\mu > 1, \epsilon > 0$

↳ • Solve

$$\min_x \quad t f(x) + q(x)$$
$$\text{st} \quad Ax - b = 0$$

(using, say ALM; ADMM; etc)

• if $\frac{m}{t} < \epsilon$ quit, else $t = \mu t$

⑧

Now look at problem

$$\min_x \quad t f(x) + \varphi(x)$$

$$\text{st } Ax = b$$

When we solve, get x^*, λ_i^*, v^*
and

$$Ax - b = 0$$

$$g_i(x) \leq 0$$

$$\lambda_i^* \geq 0$$

$$\nabla f + \sum_i \lambda_i^* \nabla g_i + v^{*\top} A = 0$$

$$-\lambda_i^* g_i = \frac{1}{t}$$

↑
here, we got λ^*, v^* from
 x^* — what if they
are vars!

KKT for original
problem

$$Ax - b = 0$$

$$g_i(x) \leq 0$$

$$\lambda_i \geq 0$$

$$\nabla f + \sum_i \lambda_i \nabla g_i + v^{\top} A = 0$$

$$\lambda_i g_i = 0$$